

# On $\mathcal{F}$ -independence in Graphs

F. Göring<sup>1</sup>

Jochen Harant<sup>2</sup>

Dieter Rautenbach<sup>2</sup>

and

Ingo Schiermeyer<sup>3</sup>

<sup>1</sup> Fakultät für Mathematik, TU Chemnitz, 09107 Chemnitz, Germany,  
email: frank.goering@mathematik.tu-chemnitz.de

<sup>2</sup> Institut für Mathematik, TU Ilmenau, Postfach 100565, 98684 Ilmenau, Germany,  
email: jochen.harant@tu-ilmenau.de, dieter.rautenbach@tu-ilmenau.de

<sup>3</sup> Institut für Diskrete Mathematik und Algebra, TU Bergakademie Freiberg, 09596  
Freiberg, Germany, email: schierme@math.tu-freiberg.de

**Abstract.** Let  $\mathcal{F}$  be a set of graphs and for a graph  $G$  let  $\alpha_{\mathcal{F}}(G)$  and  $\alpha_{\mathcal{F}}^*(G)$  denote the maximum order of an induced subgraph of  $G$  which does not contain a graph in  $\mathcal{F}$  as a subgraph and which does not contain a graph in  $\mathcal{F}$  as an induced subgraph, respectively. Lower bounds on  $\alpha_{\mathcal{F}}(G)$  and  $\alpha_{\mathcal{F}}^*(G)$  and algorithms realizing them are presented.

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## 1 Introduction

We consider finite, undirected and simple graphs  $G$  with vertex set  $V(G)$  and edge set  $E(G)$  and refer to [5] for undefined notation.

As a generalization of the well-studied concept of independent sets [8] in graphs Peter Mihok [9] proposed the following problem: *For two given graphs  $F$  and  $G$ , what is the maximum order of an induced subgraph of  $G$  that either does not contain  $F$  as a subgraph or does not contain  $F$  as an induced subgraph?*

The purpose of the present paper is to formalize the independence concept corresponding to this problem and to initiate its study. Therefore, for a graph  $G$  and a set  $\mathcal{M}$  of graphs we denoted by  $f(G, \mathcal{M})$  the maximum order  $|S|$  of a subgraph  $G[S]$  of  $G$  induced by  $S \subseteq V(G)$  such that  $G[S]$  belongs to  $\mathcal{M}$ . Choosing  $\mathcal{M}$  appropriately allows to capture Mihok's independence problem. More precisely, let  $\mathcal{F}$  be a set of graphs and for a graph  $G$  let  $\alpha_{\mathcal{F}}(G)$  and  $\alpha_{\mathcal{F}}^*(G)$  denote the maximum order of an induced subgraph of  $G$  which does not contain a graph in  $\mathcal{F}$  as a subgraph and which does not contain a graph in  $\mathcal{F}$  as an induced subgraph, respectively. Clearly, if we define  $\mathcal{M}_{\mathcal{F}}$  as the set of all graphs which do not contain a graph in  $\mathcal{F}$  as a subgraph and  $\mathcal{M}_{\mathcal{F}}^*$  as the set of all graphs which do not contain a graph in  $\mathcal{F}$  as an induced subgraph, then  $\alpha_{\mathcal{F}}(G) = f(G, \mathcal{M}_{\mathcal{F}})$  and  $\alpha_{\mathcal{F}}^*(G) = f(G, \mathcal{M}_{\mathcal{F}}^*)$ . If  $\mathcal{F} = \{F\}$ , then we write  $\alpha_F(G)$  and  $\alpha_F^*(G)$  for short.

Several well-known graph parameters are special cases of these notions as shown in the following result which collects some obvious basic observations.

**Proposition 1** *Let  $G$  be a graph.*

- (i)  $\alpha_{K_2}(G)$  equals the independence number  $\alpha(G)$  of  $G$ .
- (ii)  $\alpha_{\bar{K}_2}(G)$  equals the clique number of  $G$ .
- (iii)  $\alpha_{P_3}(G)$  equals the dissociation number of  $G$  [2].
- (iv)  $\alpha_{K_r}(G) = \alpha_{K_r}^*(G)$ .
- (v)  $\alpha_{\bar{K}_r}(G) = \min\{|V(G)|, r - 1\}$ .
- (vi)  $\alpha_{\bar{K}_r}^*(G) = \max\{|S| \mid S \subseteq V(G), \alpha(G[S]) \leq r - 1\}$ .
- (vii)  $\alpha_{\mathcal{F}}^*(G) = \alpha_{\{\bar{F} \mid F \in \mathcal{F}\}}^*(\bar{G})$ .

Our next result is a lower bound on  $f(G, \mathcal{M})$  provided the set  $\mathcal{M}$  has some natural properties.

**Theorem 2** *Let  $\mathcal{M}$  be a set of graphs and let  $G$  be a graph.*

- (i) *If  $\mathcal{M}$  is closed under taking induced subgraphs, then*

$$f(G, \mathcal{M}) \geq \sum_{S: S \subseteq V(G), G[S] \in \mathcal{M}} \binom{|V(G)|}{|S|}^{-1}$$

- (ii) *If  $\mathcal{M}$  is closed under taking induced subgraphs and under forming the union of graphs, then*

$$f(G, \mathcal{M}) \geq \sum_{S: S \subseteq V(G), G[S] \in \mathcal{M}, G[S] \text{ is connected}} \binom{|N_G[S]|}{|S|}^{-1}$$

where  $N_G[S] = \cup_{u \in S} N_G[u]$ .

*Proof:* We only prove (ii) and leave the very similar proof of (i) to the reader. We choose a permutation  $v_1, v_2, \dots, v_n$  of the vertices of  $G$  uniformly at random. Let  $S_0 = \emptyset$  and for  $1 \leq i \leq n$  let  $S_i = S_{i-1} \cup \{v_i\}$  if  $G[S_{i-1} \cup \{v_i\}] \in \mathcal{M}$  and  $S_i = S_{i-1}$  otherwise. Clearly,  $f(G, \mathcal{M}) \geq |S_n|$  and  $v_i \in S_n$  if and only if  $v_i \in S_i$  and the component  $H_i$  of  $G[S_i]$  containing  $v_i$  belongs to  $\mathcal{M}$ . Therefore, for a set  $S \subseteq V(G)$  with  $v_i \in S$  such that  $G[S] \in \mathcal{M}$  and  $G[S]$  is connected, a lower bound for the probability that  $H_i = G[S]$  is the probability that in

the chosen permutation the vertices  $S \setminus \{v_i\}$  precede  $v_i$  while  $v_i$  precedes the vertices in  $N_G[S] \setminus S$  which equals  $\frac{1}{|S|} \binom{|N_G[S]|}{|S|}^{-1}$ . Therefore, by linearity of expectation

$$\begin{aligned}
f(G, \mathcal{M}) &\geq \mathbf{E}(|S_n|) = \sum_{i=1}^n \mathbf{P}(v_i \in S_n) \\
&\geq \sum_{i=1}^n \sum_{S: v_i \in S \subseteq V(G), G[S] \in \mathcal{M}, G[S] \text{ is connected}} \frac{1}{|S|} \binom{|N_G[S]|}{|S|}^{-1} \\
&= \sum_{S: S \subseteq V(G), G[S] \in \mathcal{M}, G[S] \text{ is connected}} \sum_{i: v_i \in S} \frac{1}{|S|} \binom{|N_G[S]|}{|S|}^{-1} \\
&= \sum_{S: S \subseteq V(G), G[S] \in \mathcal{M}, G[S] \text{ is connected}} \binom{|N_G[S]|}{|S|}^{-1}
\end{aligned}$$

and the proof is complete.  $\square$

**Corollary 3** *Let  $G$  be a graph.*

(i)  $\alpha(G) \geq \sum_{u \in V(G)} \frac{1}{1+d_G(v)}$  (Caro [3], Wei [11]).

(ii) *The dissociation number satisfies*

$$\alpha_{P_3}(G) \geq \sum_{u \in V(G)} \frac{1}{1+d_G(v)} + \sum_{uv \in E(G)} \frac{2}{|N_G[u] \cup N_G[v]| (|N_G[u] \cup N_G[v]| - 1)}.$$

*Proof:* Note that  $\mathcal{M}_{\{K_2\}} = \{\bar{K}_r \mid r \in \mathbb{N}\}$  and  $\mathcal{M}_{\{P_3\}} = \mathcal{M}_{\{K_2\}} \cup \{K_2 \cup \bar{K}_r \mid r \in \mathbb{N}\}$ . Both statements follow immediately from Theorem 2(ii) and the observation that the only connected graph in  $\mathcal{M}_{\{K_2\}}$  is  $K_1$  and the only connected graphs in  $\mathcal{M}_{\{P_3\}}$  are  $K_1$  and  $K_2$ .  $\square$

The famous bound due to Caro [3] and Wei [11] from Corollary 3 has yet another generalization in this context.

**Theorem 4** *If  $G$  is a graph and  $r \in \mathbb{N}$ , then  $\alpha_{K_{r+1}}(G) \geq \sum_{v \in V(G)} \frac{1}{1+d_G(v) - \alpha_{K_r}(G[N_G(v)])}$ .*

*Proof:* We mimic a proof from [1]. For every vertex  $v \in V(G)$  let the set  $X_v \subseteq N_G(v)$  be such that  $|X_v| = d_G(v) - \alpha_{K_r}(G[N_G(v)])$  and  $G[N_G(v) \setminus X_v]$  does not contain  $K_r$  as a subgraph. Let  $v_1, v_2, \dots, v_n$  be a permutation of the vertices of  $G$  chosen uniformly at random and let  $v_i \in S$  if and only if  $X_{v_i} \cap \{v_1, v_2, \dots, v_i\} = \emptyset$ , i.e.  $v_i$  is the first vertex of  $\{v_i\} \cup X_{v_i}$  that appears within the permutation. Clearly,  $G[S]$  does not contain  $K_{r+1}$  as a subgraph and

$$\alpha_{K_{r+1}}(G) \geq \mathbf{E}(|S|) = \sum_{v \in V(G)} \mathbf{P}(v \in S) = \sum_{v \in V(G)} \frac{1}{1+d_G(v) - \alpha_{K_r}(G[N_G(v)])}. \square$$

The next result relies on methods proposed in [7].

**Theorem 5** *If  $G$  is a graph with vertex set  $\{v_1, v_2, \dots, v_n\}$  and  $r \in \mathbb{N}$ , then*

$$\alpha_{K_{1,r}}(G) = \max_{v_i \in V(G)} \sum_{Y: Y \subseteq N_G(v_i), |Y| < r} p_i \left( \prod_{v_j \in Y} p_j \prod_{v_k \in N_G(v_i) \setminus Y} (1 - p_j) \right),$$

where the maximum is taken over all  $(p_1, p_2, \dots, p_n) \in [0, 1]^n$ .

*Proof:* Let  $p_i \in [0, 1]$  for  $1 \leq i \leq n$ . We consider a random subset  $X$  of  $V(G)$  formed by choosing every vertex  $v_i$  independently with probability  $p_i$ . If  $S = \{v \in X \mid d_{G[X]}(v) < r\}$ , then

$$\alpha_{K_{1,r}}(G) \geq \mathbf{E}(S) = \sum_{v_i \in V(G)} p_i \sum_{Y: Y \subseteq N_G(v_i), |Y| < r} \left( \prod_{v_j \in Y} p_j \prod_{v_k \in N_G(v_i) \setminus Y} (1 - p_j) \right).$$

Conversely, if  $S \subseteq V$  is such that  $\alpha_{K_{1,r}}(G) = |S|$  and  $G[S]$  has maximum degree less than  $r$ , then setting  $p_i^* = 1$  for all  $v_i \in S$  and  $p_i^* = 0$  for all  $v_i \notin S$  yields

$$\alpha_{K_{1,r}}(G) = \mathbf{E}(S) = \sum_{v_i \in V(G)} p_i^* \sum_{Y: Y \subseteq N_G(v_i), |Y| < r} \left( \prod_{v_j \in Y} p_j^* \prod_{v_k \in N_G(v_i) \setminus Y} (1 - p_j^*) \right)$$

which completes the proof.  $\square$

It is trivial that for several specific choices of  $\mathcal{M}$  and  $\mathcal{F}$  the decision problems associated with  $f(G, \mathcal{M})$ ,  $\alpha_{\mathcal{F}}(G)$  and  $\alpha_{\mathcal{F}}^*(G)$  are NP-complete. In view of Mihok's original problem, we consider the case that  $\mathcal{F}$  consists of just one graph in more detail.

**Theorem 6** *If  $F$  is a graph containing at least one edge, then the following problems are NP-complete problem.*

- (i) *For a given graph  $G$  and  $k \in \mathbb{N}$ , decide whether  $\alpha_F(G) \geq k$ .*
- (ii) *For a given graph  $G$  and  $k \in \mathbb{N}$ , decide whether  $\alpha_F^*(G) \geq k$ .*

*Proof:* Let  $uv$  be an arbitrary edge of  $F$ . For a graph  $G$  let the graph  $G'$  arise as follows: For every edge  $xy$  of  $G$  add a copy  $F_{xy}$  of  $F$  and identify the copy of the edge  $uv$  in  $F_{xy}$  with  $xy$  (in any orientation).

It is obvious that for every set  $T \subseteq V(G')$  of minimum cardinality such that  $G'[V(G') \setminus T]$  does not contain  $F$  as a subgraph (or induced subgraph),  $T$  must intersect every copy  $F_{xy}$  of  $F$  in  $G'$ . Since deleting either  $x$  or  $y$  from  $F_{xy}$  clearly deletes this copy of  $F$ , we can assume that  $T \subseteq V(G)$  and that  $T \cap \{x, y\} \neq \emptyset$  for all  $xy \in E(G)$ . Hence  $T$  is exactly a vertex cover of  $G$ . This implies  $\alpha(G) = \alpha_F(G') = \alpha_F^*(G')$  and the desired statement follows from

the NP-completeness of the corresponding decision problem for the independence number [6].  $\square$

Note that in view Proposition 1(vii), the decision problem “ $\alpha_{\mathcal{F}}^*(G) \geq k$ ?” remains NP-complete even if  $F$  is edge-less.

Tuza [10] observed the following nice relation between the independence number and the domination number  $\gamma(G)$  of a graph  $G$  [7]:

$$\alpha(G) = \max\{\gamma(H) \mid H \text{ is an induced subgraph of } G\}.$$

We close with a generalization of this equality. For a set  $\mathcal{F}$  of graphs and a graph  $G$  let  $\gamma_{\mathcal{F}}(G)$  ( $\gamma_{\mathcal{F}}^*(G)$ ) denote the minimum cardinality  $|D|$  of a set  $D \subseteq V(G)$  such that for every vertex  $u \in V(G) \setminus D$  there is a graph  $F \in \mathcal{F}$  and a set  $D' \subseteq D$  with  $|D'| = |V(F)| - 1$  such that  $G[D' \cup \{u\}]$  contains a graph in  $\mathcal{F}$  as a(n induced) subgraph.

**Theorem 7** *If  $\mathcal{F}$  is a set of graphs and let  $G$  is a graph  $G$ , then*

$$\begin{aligned} \alpha_{\mathcal{F}}(G) &= \max\{\gamma_{\mathcal{F}}(H) \mid H \text{ is an induced subgraph of } G\} \\ \alpha_{\mathcal{F}}^*(G) &= \max\{\gamma_{\mathcal{F}}^*(H) \mid H \text{ is an induced subgraph of } G\}. \end{aligned}$$

*Proof:* We only prove the first equality and leave the very similar proof of the second equality to the reader.

If  $S \subseteq V(G)$  is such that  $|S| = \alpha_{\mathcal{F}}(G)$  and  $G[S]$  does not contain a graph in  $\mathcal{F}$  as a subgraph, then  $\gamma_{\mathcal{F}}(G[S]) = |S| \geq \alpha_{\mathcal{F}}(G)$ .

Conversely, if  $G[S]$  is an induced subgraph of  $G$  for which  $\gamma_{\mathcal{F}}(G[S])$  is maximum, then let  $S' \subseteq S$  be of maximum cardinality such that  $G[S']$  does not contain a graph in  $\mathcal{F}$  as a subgraph. We obtain  $\gamma_{\mathcal{F}}(G[S]) \leq |S'| = \alpha_{\mathcal{F}}(G[S]) \leq \alpha_{\mathcal{F}}(G)$  and the proof is complete.  $\square$

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