

Interpolating between Bounds on the Independence Number

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Abstract

For a non-negative integer T , we prove that the independence number of a graph $G = (V, E)$ in which every vertex belongs to at most T triangles is at least $\sum_{u \in V} f(d(u), T)$ where $d(u)$ denotes the degree of a vertex $u \in V$, $f(d, T) = \frac{1}{d+1}$ for $T \geq \binom{d}{2}$ and $f(d, T) = (1 + (d^2 - d - 2T)f(d-1, T))/(d^2 + 1 - 2T)$ for $T < \binom{d}{2}$. This is a common generalization of the lower bounds for the independence number due to Caro, Wei, and Shearer. We discuss further possible strengthenings of our result and pose a corresponding conjecture.

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1 Introduction

We consider finite, simple, and undirected graphs $G = (V, E)$ with vertex set V and edge set E . The degree of a vertex u in G is denoted by $d_G(u)$. A set of vertices $I \subseteq V$ of G is called independent, if no two vertices in I are adjacent. The independence number $\alpha(G)$ is the maximum cardinality of an independent set.

The independence number is among the most fundamental and well-studied graph-theoretical concepts. In view of its computational hardness [7] bounds on the independence number received a lot of attention. The following classical lower bound on the independence number of a graph G was obtained independently by Caro [4] and Wei [13]

$$\alpha(G) \geq \sum_{u \in V} \frac{1}{d_G(u) + 1}. \quad (1)$$

This bound is best-possible in view of cliques. A simple proof of (1) is based on the observation that the deletion of a vertex of maximum degree at least 1 from G does not decrease the right-hand side of (1). Therefore, iteratively deleting such vertices results in an independent set of at least the desired cardinality.

For triangle-free graphs G , Shearer [11] (cf. also [10]) proved

$$\alpha(G) \geq \sum_{u \in V} f(d_G(u)) \tag{2}$$

where $f(0) = 1$ and $f(d) = \frac{1+(d^2-d)f(d-1)}{d^2+1}$ for $d \in \mathbb{N}$. The bound (2) improved on earlier results [2, 3, 6] which gave bounds of the form $\alpha(G) \geq \Omega\left(\frac{n \ln(d)}{d}\right)$ for triangle-free graph G of order n and average degree d . For related results concerning k -clique-free graphs, we refer to [1, 9, 12].

Shearer's bound (2) is similar to Caro and Wei's bound (1) in the sense that every vertex contributes a suitable degree-dependent weight to the value of the bound. Its inductive proof is considerably harder than the proof for (1). In [11] Shearer exploited his approach further to establish lower bounds on the independence number of graphs of large girth. For d -regular graphs G of order n and girth g , he proved $\alpha(G) \geq (1 - o(g))nf(d)$ where $f(3) = \frac{125}{302}$ and $f(d) = \frac{1+(d^2-d)f(d-1)}{d^2+1}$ for $d \geq 4$. The strength of his approach is illustrated by the fact that this last bound was only improved very recently [5, 8].

The goal of the research reported here was to prove a common generalization of (1) and (2). For a graph G and a vertex u of G , let $t_G(u)$ denote the number of triangles of G containing u . Note that $t_G(u)$ equals the number of edges among neighbours of u in G . For a suitable function $f : \mathbb{N}_0^2 \rightarrow \mathbb{R}_{\geq 0}$, we wanted to prove a bound of the form

$$\alpha(G) \geq \sum_{u \in V} f(d_G(u), t_G(u))$$

which coincides with (2) for triangle-free graphs and is always at least as good as (1).

In Section 2 we discuss Shearer's approach and the possibility to extend it to graphs which may contain triangles. This leads to a number of properties the function f should possess. In Section 3 we propose a candidate for f and establish most of the desired properties. While we eventually succeed in proving a common generalization of (1) and (2), we found our result not yet totally satisfactory and pose a conjecture concerning a possible strengthening.

2 Extending Shearer's Approach

In this section we discuss how to extend Shearer's approach from [11] to graphs which may contain triangles. Consider a graph G . For a vertex u in G , let $d_u = d_G(u)$ and $t_u = t_G(u)$. Our goal is a lower bound for the independence number of G of the form

$$\alpha(G) \geq w(G) := \sum_{v \in V} f(d_v, t_v) \tag{3}$$

where $f : \mathbb{N}_0^2 \rightarrow \mathbb{R}_{\geq 0}$ is a suitable function. In order for Shearer's inductive approach to work, the function f has to possess several properties. For $d, t \in \mathbb{N}_0$, we assume

$$(P_1) \quad f(0, 0) = 1,$$

$$(P_2) \quad f(d, t) \geq f(d, t + 1),$$

$$(P_3) \quad f(d, t) - f(d + 1, t) \geq f(d + 1, t) - f(d + 2, t), \text{ and}$$

(P_4) $1 - (d+2)f(d+1, t) + ((d+1)^2 - (d+1) - 2t)(f(d, t) - f(d+1, t)) \geq 0$ for $t \leq \binom{d+1}{2}$.

Property (P_1) implies (3) for $|V| = 1$, i.e. the base case of the induction. Furthermore, by (P_1), we may assume that G has no vertex of degree 0.

For two distinct vertices u and v in G , let $d_{\{u,v\}}$ denote the number of common neighbours of u and v . For a vertex u in G , let N_u denote the set of neighbours of u and let N_u^2 denote the set of vertices at distance exactly two from u , respectively.

If there is a vertex u in G such that the deletion of all vertices in $\{u\} \cup N_u$ results in a graph G_u with $1 - w(G) + w(G_u) \geq 0$, then adding u to a maximum independent set of G_u results in an independent set of G of order at least $1 + w(G_u) \geq w(G)$. If $w \in N_u^2$, then $d_{G_u}(w) = d_w - d_{\{u,w\}}$ and $t_{G_u}(w) \leq t_w$. Therefore, by the monotonicity property (P_2), it suffices to prove the existence of a vertex u in G with

$$1 - f(d_u, t_u) - \sum_{v \in N_u} f(d_v, t_v) + \sum_{w \in N_u^2} (f(d_w - d_{\{u,w\}}, t_w) - f(d_w, t_w)) \geq 0. \quad (4)$$

In [11] Shearer shows the existence of such a vertex by proving that (4) holds on average. Therefore, let

$$A = \sum_{u \in V} \left(1 - f(d_u, t_u) - \sum_{v \in N_u} f(d_v, t_v) + \sum_{w \in N_u^2} (f(d_w - d_{\{u,w\}}, t_w) - f(d_w, t_w)) \right).$$

Since $\sum_{u \in V} \sum_{v \in N_u} f(d_v, t_v) = \sum_{u \in V} d_u f(d_u, t_u)$ and $w \in N_u^2 \Leftrightarrow u \in N_w^2$, we have

$$\begin{aligned} A &= \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, t_u) + \sum_{w \in N_u^2} (f(d_w - d_{\{u,w\}}, t_w) - f(d_w, t_w)) \right) \\ &= \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, t_u) + \sum_{w \in N_u^2} (f(d_u - d_{\{u,w\}}, t_u) - f(d_u, t_u)) \right). \end{aligned} \quad (5)$$

By (P_3),

$$f(d_u - d_{\{u,w\}}, t_u) - f(d_u, t_u) \geq d_{\{u,w\}}(f(d_u - 1, t_u) - f(d_u, t_u)).$$

Furthermore, simple double-counting yields

$$\sum_{w \in N_u^2} d_{\{u,w\}} = \left(\sum_{v \in N_u} (d_v - 1) \right) - 2t_u.$$

Together with (5) we obtain

$$\begin{aligned} A &\geq \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, t_u) + \sum_{w \in N_u^2} d_{\{u,w\}}(f(d_u - 1, t_u) - f(d_u, t_u)) \right) \\ &= \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, t_u) + \left(\left(\sum_{v \in N_u} (d_v - 1) \right) - 2t_u \right) (f(d_u - 1, t_u) - f(d_u, t_u)) \right). \end{aligned} \quad (6)$$

A crucial property of f — or of the pair (G, f) — needed at this point to continue along Shearer's argument is that

$$\sum_{u \in V} \sum_{v \in N_u} (d_v - 1)(f(d_u - 1, t_u) - f(d_u, t_u)) \geq \sum_{u \in V} \sum_{v \in N_u} (d_u - 1)(f(d_u - 1, t_u) - f(d_u, t_u)). \quad (7)$$

If the values of f are independent of the second parameter, i.e. $f(d, t) = f(d, t + 1)$ for all $d, t \in \mathbb{N}_0$, then (7) follows from property (P_3) as follows

$$\begin{aligned} & \sum_{u \in V} \sum_{v \in N_u} (d_v - 1)(f(d_u - 1, t_u) - f(d_u, t_u)) \\ &= \sum_{uv \in E} ((d_v - 1)(f(d_u - 1, t_u) - f(d_u, t_u)) + (d_u - 1)(f(d_v - 1, t_v) - f(d_v, t_v))) \\ &\stackrel{(P_3)}{\geq} \sum_{uv \in E} ((d_u - 1)(f(d_u - 1, t_u) - f(d_u, t_u)) + (d_v - 1)(f(d_v - 1, t_v) - f(d_v, t_v))) \\ &= \sum_{u \in V} \sum_{v \in N_u} (d_u - 1)(f(d_u - 1, t_u) - f(d_u, t_u)). \end{aligned}$$

Assuming (7) we would obtain from (6) that

$$\begin{aligned} A &\geq \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, t_u) + \left(\left(\sum_{v \in N_u} (d_v - 1) \right) - 2t_u \right) (f(d_u - 1, t_u) - f(d_u, t_u)) \right) \\ &= \sum_{u \in V} (1 - (d_u + 1)f(d_u, t_u) + (d_u^2 - d_u - 2t_u) (f(d_u - 1, t_u) - f(d_u, t_u))). \end{aligned}$$

Since $t_u \leq \binom{d_u}{2}$ for every vertex u in G , property (P_4) would imply $A \geq 0$ which would complete the inductive proof. In order to turn the sketched approach into a result we need to describe a function f which possesses the desired properties. In fact, apart from a version of (7) in full generality our proposal for f will possess all these properties.

3 A Reasonable Proposal for f

In this section we propose a choice for f which has properties (P_1) through (P_4) and which appears reasonable in the sense that it allows to prove a common generalization of Caro and Wei's bound (1) and Shearer's bound (2).

For non-negative integers d and t let

$$r(d, t, f) = \frac{1 + (d^2 - d - 2t)f}{d^2 + 1 - 2t}. \quad (8)$$

Furthermore, let

$$f(d, t) = \begin{cases} \frac{1}{\binom{d+1}{2}} & , t \geq \binom{d}{2}, \\ r(d, t, f(d-1, t)) & , t < \binom{d}{2}. \end{cases} \quad (9)$$

Clearly, the function $f(\cdot, 0)$ coincides with the function $f(\cdot)$ from (2). Furthermore, we will show $f(d, t) \geq \frac{1}{d+1}$ for $d, t \in \mathbb{N}_0$. In view of Section 2 it makes sense to define $f(d, t)$ also for values of d

and t with $t > \binom{d}{2}$ which are graph-theoretically meaningless. Table 1 shows some specific values of f . The bold entries correspond to vertices whose neighbourhoods induce complete graphs. As soon as the neighbourhood of a vertex is not complete the Shearer-like recursion (8) sets in.

$f(d, t)$	$d = 0$	$d = 1$	$d = 2$	$d = 3$	$d = 4$	$d = 5$
$t = 0$	1	1/2	2/5	17/50	127/425	593/2210
$t = 1$	1	1/2	1/3	7/24	47/180	19/80
$t = 2$	1	1/2	1/3	5/18	29/117	581/2574
$t = 3$	1	1/2	1/3	1/4	5/22	23/110
$t = 4$	1	1/2	1/3	1/4	2/9	11/54
$t = 5$	1	1/2	1/3	1/4	3/14	11/56
$t = 6$	1	1/2	1/3	1/4	1/5	13/70
$t = 7$	1	1/2	1/3	1/4	1/5	11/60
$t = 8$	1	1/2	1/3	1/4	1/5	9/50
$t = 9$	1	1/2	1/3	1/4	1/5	7/40
$t = 10$	1	1/2	1/3	1/4	1/5	1/6

Table 1 $f(d, t)$ for $0 \leq d \leq 5$ and $0 \leq t \leq 10$.

The next lemma collects properties of f . For $t \in \mathbb{N}_0$, let $d_t = \max \left\{ d \in \mathbb{N}_0 \mid t \geq \binom{d}{2} \right\}$. Note that (9) is equivalent with $f(d, t) = \frac{1}{d+1}$ for $d \leq d_t$ and $f(d, t) = r(d, t, f(d-1, t))$ for $d > d_t$.

Lemma 1 *Let $d, t \in \mathbb{N}_0$.*

(i) $f(d, t) \geq \frac{3(d+2)}{2(d^2+5d+5+t)}$ for $d \geq d_t$.

(ii) $f(d, t) \geq \frac{1}{d+1}$.

(iii) $f(d, t) \geq f(d+1, t)$.

(iv) $f(d, t) \geq f(d, t+1)$.

(v) $f(d, t) - f(d+1, t) \geq f(d+1, t) - f(d+2, t)$.

(vi) $1 - (d+2)f(d+1, t) + ((d+1)^2 - (d+1) - 2t)(f(d, t) - f(d+1, t)) \geq 0$ for $t \leq \binom{d+1}{2}$.

Proof: (i) We prove this statement by induction on $d \geq d_t$. By (9), $f(d_t, t) = \frac{1}{d_t+1}$. Since $t \geq \binom{d_t}{2} = \frac{1}{2}d_t(d_t - 1)$, we obtain

$$\begin{aligned} \frac{3(d_t+2)}{2(d_t^2+5d_t+5+t)} &\leq \frac{3(d_t+2)}{2(d_t^2+5d_t+5+\frac{1}{2}d_t(d_t-1))} \\ &= \frac{3(d_t+2)}{3(d_t+2)(d_t+1)+4} < \frac{1}{d_t+1} = f(d_t, t) \end{aligned}$$

which proves the base case of the induction.

If $d > d_t$, then $2t < d(d-1)$ and, by (9), $f(d, t) = r(d, t, f(d-1, t))$. Since $r(d, t, x)$ is monotonously decreasing as a function of x , we obtain, by induction,

$$\begin{aligned}
f(d, t) - \frac{3(d+2)}{2(d^2 + 5d + 5 + t)} &= r(d, t, f(d-1, t)) - \frac{3(d+2)}{2(d^2 + 5d + 5 + t)} \\
&\geq r\left(d, t, \frac{3((d-1)+2)}{2((d-1)^2 + 5(d-1) + 5 + t)}\right) - \frac{3(d+2)}{2(d^2 + 5d + 5 + t)} \\
&= \frac{d^2 + 2d + 2 - 2t}{(d^2 + 5d + 5 + t)(d^3 + 3d + 1 + t)} \\
&\geq \frac{d^2 + 2d + 2 - d(d-1)}{(d^2 + 5d + 5 + t)(d^3 + 3d + 1 + t)} \\
&= \frac{3d+2}{(d^2 + 5d + 5 + t)(d^3 + 3d + 1 + t)} > 0
\end{aligned}$$

which completes the proof of (i).

(ii) We prove this statement by induction on d . If $d \leq d_t$, then, by (9), $f(d, t) = \frac{1}{d+1}$.

If $d > d_t$, then $t < \binom{d}{2}$ and, by induction,

$$\begin{aligned}
f(d, t) - \frac{1}{d+1} &\stackrel{(9)}{=} r(d, t, f(d-1, t)) - \frac{1}{d+1} = \frac{1 + (d^2 - d - 2t)f(d-1, t)}{d^2 + 1 - 2t} - \frac{1}{d+1} \\
&\geq \frac{1 + (d^2 - d - 2t)\frac{1}{d}}{d^2 + 1 - 2t} - \frac{1}{d+1} = \frac{d^2 - d - 2t}{(d^2 + 1 - 2t)d(d+1)} \geq 0
\end{aligned}$$

which completes the proof of (ii).

(iii) If $d \leq d_t - 1$, then $f(d, t) = \frac{1}{d+1} > \frac{1}{d+2} = f(d+1, t)$.

If $d \geq d_t$, then $t < \binom{d+1}{2}$ and

$$f(d, t) - f(d+1, t) \stackrel{(9)}{=} f(d, t) - r(d+1, t, f(d, t)) \stackrel{(8)}{=} \frac{(d+2)f(d, t) - 1}{(d+1)^2 + 1 - 2t} \stackrel{(ii)}{\geq} 0$$

which completes the proof of (iii).

(iv) We prove this statement by induction on d . If $d \leq d_t$, then $f(d, t) \stackrel{(9)}{=} f(d, t+1) \stackrel{(9)}{=} \frac{1}{d+1}$.

Hence, we may assume that $d > d_t$ which implies $t < \binom{d}{2}$ and $f(d, t) = r(d, t, f(d-1, t))$.

If $t+1 < \binom{d}{2}$, then, by induction,

$$\begin{aligned}
f(d, t) - f(d, t+1) &\stackrel{(9)}{=} r(d, t, f(d-1, t)) - r(d, t+1, f(d-1, t+1)) \\
&= \frac{1 + (d^2 - d - 2t)f(d-1, t)}{d^2 + 1 - 2t} - \frac{1 + (d^2 - d - 2t - 2)f(d-1, t+1)}{d^2 + 1 - 2t - 2} \\
&\geq \frac{1 + (d^2 - d - 2t)f(d-1, t)}{d^2 + 1 - 2t} - \frac{1 + (d^2 - d - 2t - 2)f(d-1, t)}{d^2 + 1 - 2t - 2} \\
&= \frac{(2d+2)f(d-1, t) - 2}{(d^2 + 1 - 2t)(d^2 + 1 - 2t - 2)} \stackrel{(ii)}{\geq} 0.
\end{aligned}$$

Hence, we may assume that $t + 1 = \binom{d}{2}$. This implies

$$\begin{aligned} f(d, t) - f(d, t + 1) &\stackrel{(9)}{=} r(d, t, f(d - 1, t)) - \frac{1}{d + 1} = r\left(d, t, \frac{1}{d}\right) - \frac{1}{d + 1} \\ &= \frac{1 + (d^2 - d - 2t)\frac{1}{d}}{d^2 + 1 - 2t} - \frac{1}{d + 1} = \frac{d^2 - d - 2t}{(d^2 + 1 - 2t)d(d + 1)} \geq 0 \end{aligned}$$

which completes the proof of (iv).

(v) If $d \leq d_t - 1$, then

$$\begin{aligned} &(f(d, t) - f(d + 1, t)) - (f(d + 1, t) - f(d + 2, t)) \\ &\stackrel{(9)}{=} \left(\frac{1}{d + 1} - \frac{1}{d + 2}\right) - \left(\frac{1}{d + 2} - f(d + 2, t)\right) = f(d + 2, t) - \frac{d}{(d + 1)(d + 2)} \\ &\stackrel{(ii)}{\geq} \frac{1}{d + 3} - \frac{d}{(d + 1)(d + 2)} = \frac{2}{(d + 1)(d + 2)(d + 3)} > 0. \end{aligned}$$

If $d \geq d_t$, then

$$\begin{aligned} &(f(d, t) - f(d + 1, t)) - (f(d + 1, t) - f(d + 2, t)) \\ &\stackrel{(9)}{=} f(d, t) - 2r(d + 1, t, f(d, t)) + r(d + 2, t, r(d + 1, t, f(d, t))). \end{aligned}$$

It is straightforward to verify that the last expression is non-negative if and only if $f(d, t) \geq \frac{3(d+2)}{2(d^2+5d+5+t)}$ which holds by (i) which completes the proof of (v).

(vi) It is straightforward to verify that the desired statement is equivalent to

$$f(d + 1, t) \leq r(d + 1, t, f(d, t))$$

for $t \leq \binom{d+1}{2}$.

If $t < \binom{d+1}{2}$, this follows immediately from (9). Hence, we may assume that $t = \binom{d+1}{2}$. This implies

$$\begin{aligned} &r(d + 1, t, f(d, t)) - f(d + 1, t) \\ &= r\left(d + 1, \binom{d + 1}{2}, f\left(d, \binom{d + 1}{2}\right)\right) - f\left(d + 1, \binom{d + 1}{2}\right) \\ &= \frac{1 + \left((d + 1)^2 - (d + 1) - 2\binom{d+1}{2}\right)\frac{1}{d+1}}{(d + 1)^2 + 1 - 2\binom{d+1}{2}} - \frac{1}{d + 2} = 0 \end{aligned}$$

which completes the proof of (vi). \square

Having collected numerous properties of f we can now state a joint generalization of (1) and (2).

Theorem 2 *Let $T \in \mathbb{N}_0$. If G is a graph such that every vertex of G belongs to at most T triangles, then*

$$\alpha(G) \geq \sum_{u \in V} f(d_G(u), T).$$

Proof: We proceed by induction on the order of G as in Section 2. By Lemma 1, the function $g : \mathbb{N}_0^2 \rightarrow \mathbb{R}_{\geq 0}$ with $g(d, t) = f(d, T)$ has properties (P_1) , (P_2) , and (P_3) . Therefore, we can argue exactly as in Section 2 until the point when (6) is established (with $g(d, t) = f(d, T)$ instead of $f(d, t)$). Also as shown in Section 2, (P_3) for g implies

$$\sum_{u \in V} \sum_{v \in N_u} (d_v - 1)(f(d_u - 1, T) - f(d_u - 1, T)) \geq \sum_{u \in V} \sum_{v \in N_u} (d_u - 1)(f(d_u - 1, T) - f(d_u - 1, T)). \quad (10)$$

Starting with (6) we obtain

$$\begin{aligned} A &\geq \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, T) + \left(\left(\sum_{v \in N_u} (d_v - 1) \right) - 2t_u \right) (f(d_u - 1, T) - f(d_u, T)) \right) \\ &\stackrel{(10)}{\geq} \sum_{u \in V} \left(1 - (d_u + 1)f(d_u, T) + \left(\left(\sum_{v \in N_u} (d_u - 1) \right) - 2t_u \right) (f(d_u - 1, T) - f(d_u, T)) \right) \\ &= \sum_{u \in V} (1 - (d_u + 1)f(d_u, T) + (d_u^2 - d_u - 2t_u) (f(d_u - 1, T) - f(d_u, T))). \end{aligned}$$

If $T > \binom{d_u}{2}$, then $t_u \leq \binom{d_u}{2}$ and, by Lemma 1,

$$1 - (d_u + 1)f(d_u, T) + (d_u^2 - d_u - 2t_u) (f(d_u - 1, T) - f(d_u, T)) \geq 1 - (d_u + 1)f(d_u, T) \stackrel{(9)}{=} 0.$$

If $T \leq \binom{d_u}{2}$, then $t_u \leq T$ and, by Lemma 1,

$$\begin{aligned} &1 - (d_u + 1)f(d_u, T) + (d_u^2 - d_u - 2t_u) (f(d_u - 1, T) - f(d_u, T)) \\ &\geq 1 - (d_u + 1)f(d_u, T) + (d_u^2 - d_u - 2T) (f(d_u - 1, T) - f(d_u, T)) \geq 0. \end{aligned}$$

Altogether, we obtain $A \geq 0$ which completes the proof of the theorem. \square

Lemma 1 collected more properties than we actually needed for the proof of Theorem 2. We hope that these are helpful to prove — rather than to disprove — the following conjecture.

Conjecture 3 *If G is a graph, then $\alpha(G) \geq \sum_{u \in V} f(d_G(u), t_G(u))$.*

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